



Applied!

# Data & Network Security

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[ans.dailysec.ir](https://ans.dailysec.ir)

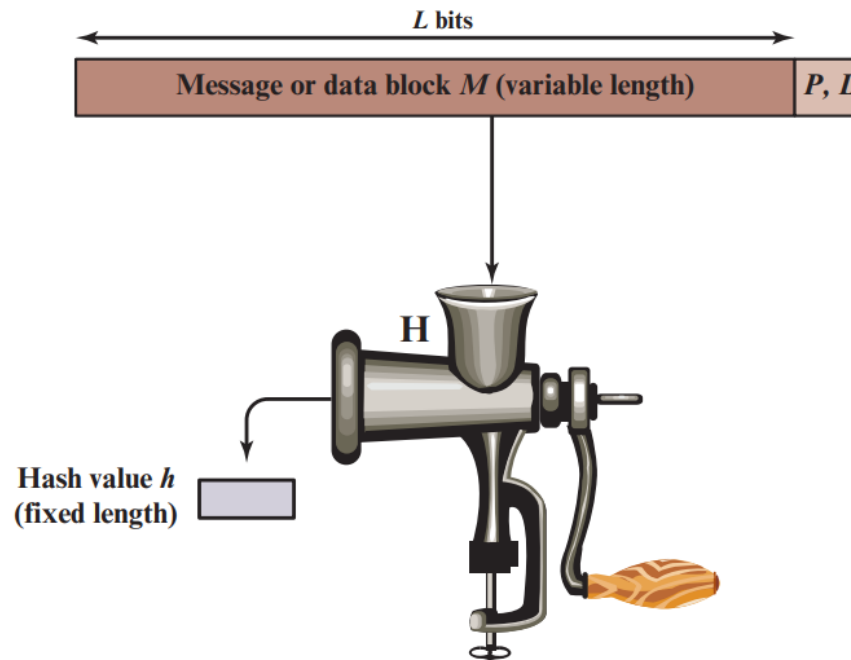
[aNetSec.github.io](https://aNetSec.github.io)

Spring 2025

# Hash Functions

# Hash Function

- A hash function  $H$  accepts a **variable-length** block of data  $M$  as input and produces a **fixed-size result**  $h = H(M)$

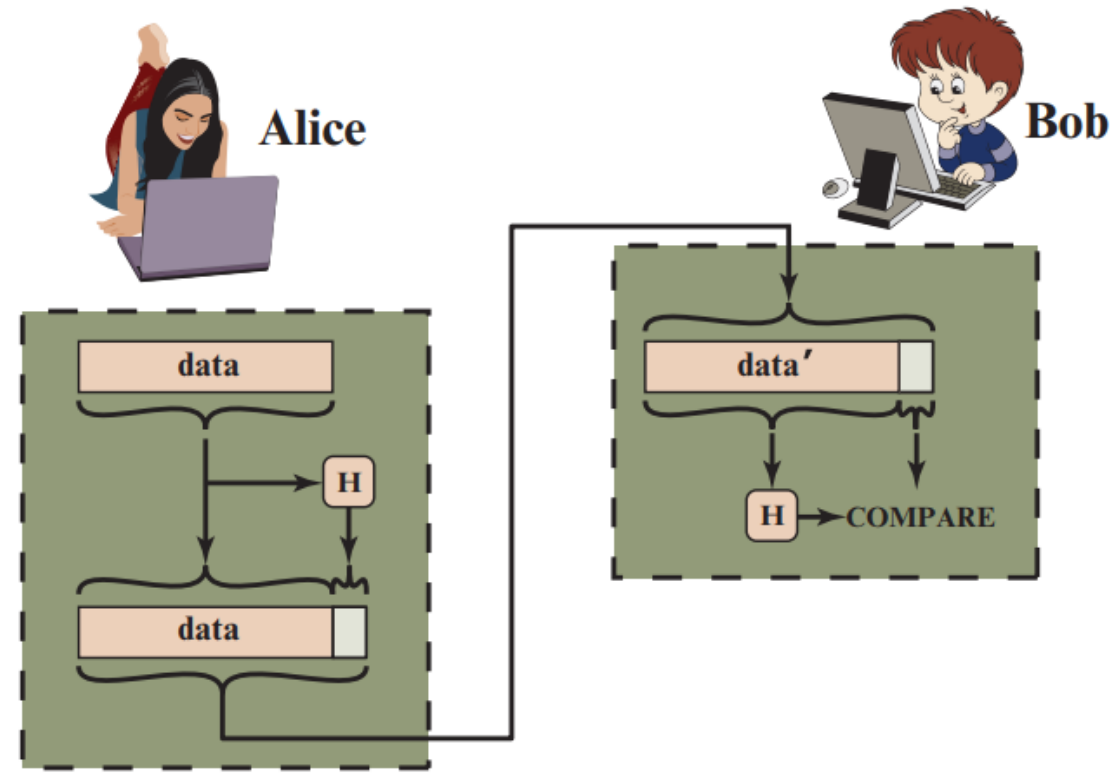


$P, L$  = padding plus length field

**Figure 11.1** Cryptographic Hash Function;  $h = H(M)$

# Applications Of Hash Functions

- Check integrity of message
- Check message not changed in transfer



(a) Use of hash function to check data integrity

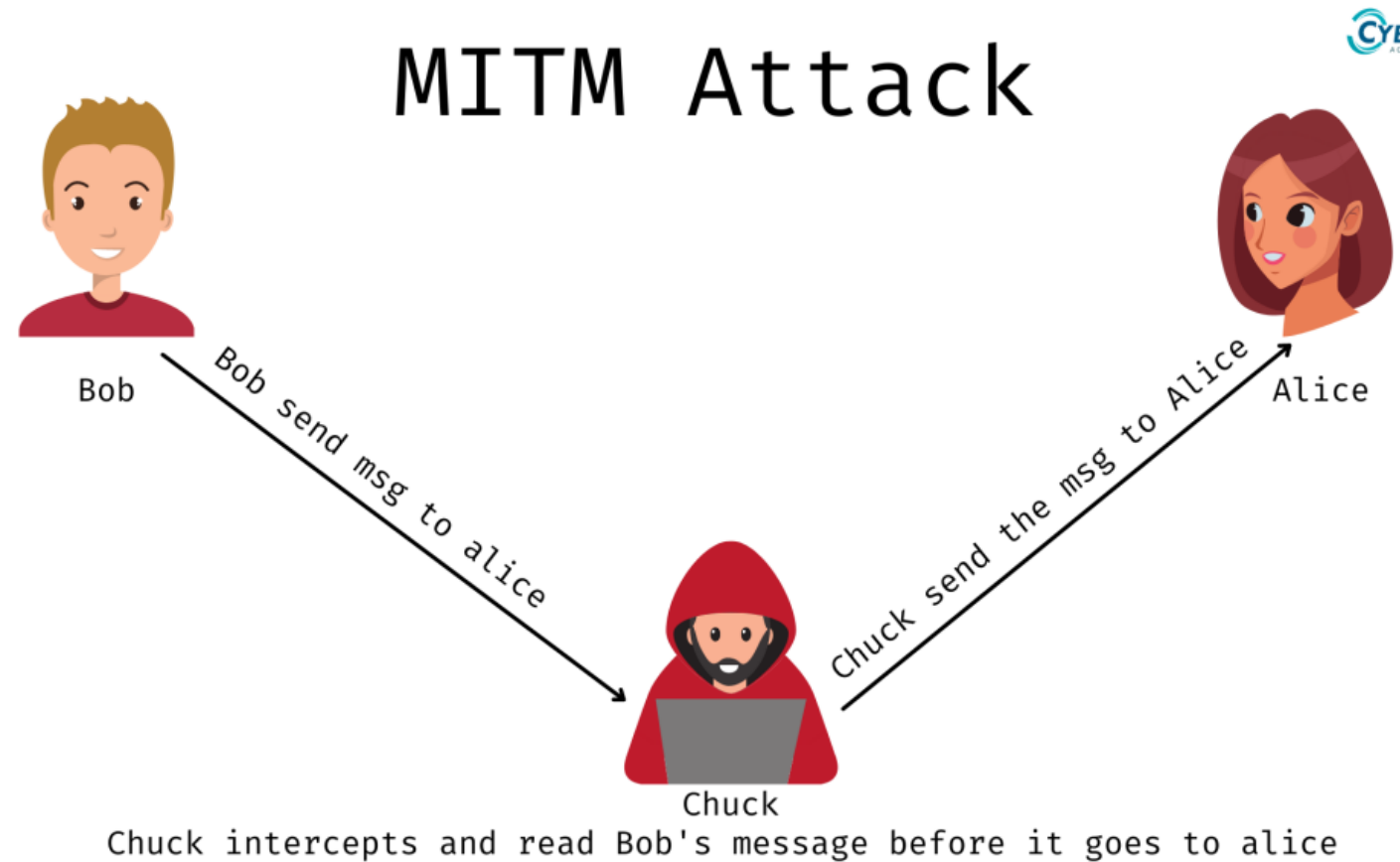
# Sniff

- Just listen to message
- No change!
- It's a passive attack

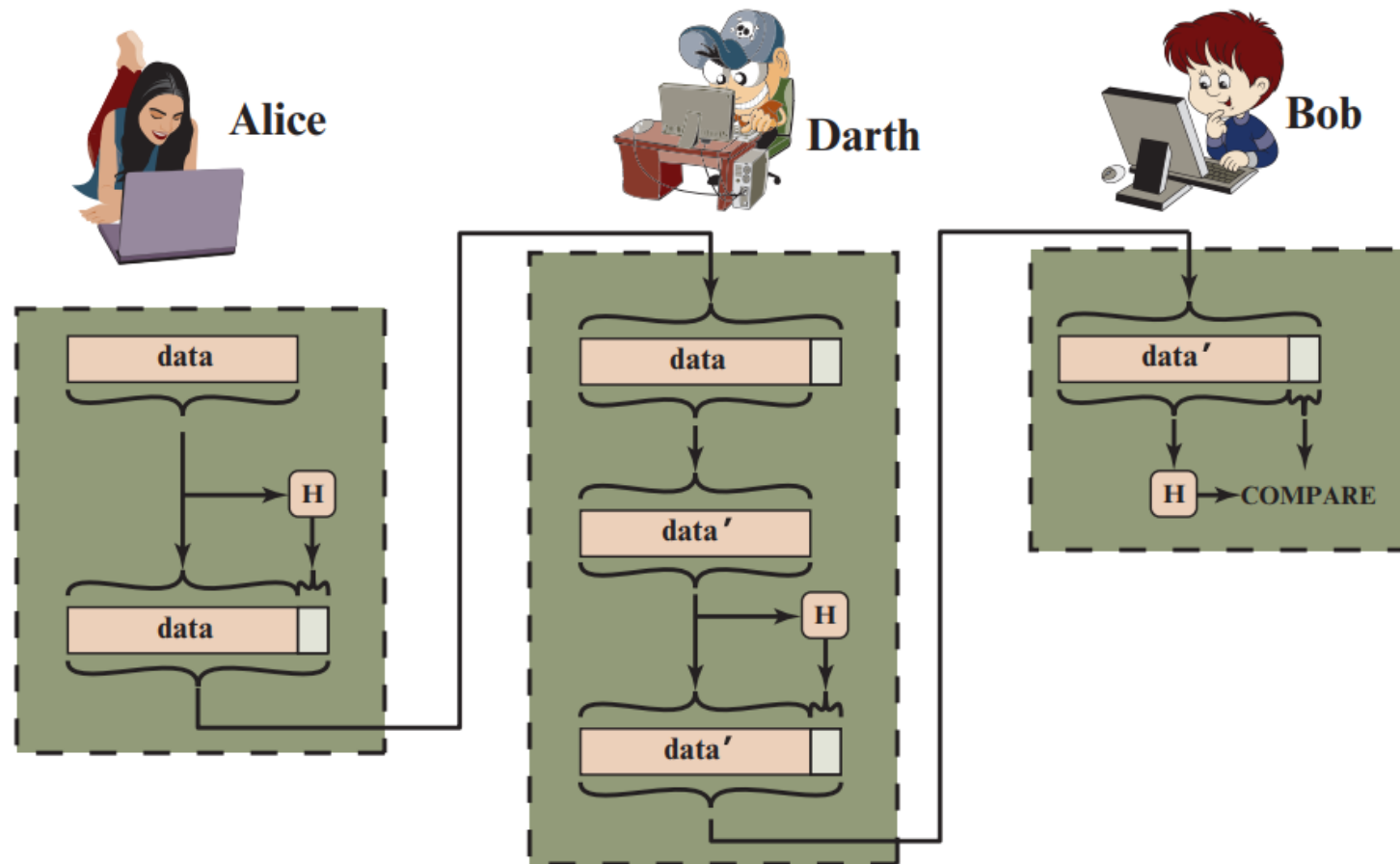


# Man-in-the-middle attack

- Change the original message
- It's an active attack



# Attack Against Hash Function



(b) Man-in-the-middle attack

# Requirements for a Cryptographic Hash Function $H$

Requirement	Description
Variable input size	$H$ can be applied to a block of data of any size.
Fixed output size	$H$ produces a fixed-length output.
Efficiency	$H(x)$ is relatively easy to compute for any given $x$ , making both hardware and software implementations practical.
Preimage resistant (one-way property)	For any given hash value $h$ , it is computationally infeasible to find $y$ such that $H(y) = h$ .
Second preimage resistant (weak collision resistant)	For any given block $x$ , it is computationally infeasible to find $y \neq x$ with $H(y) = H(x)$ .
Collision resistant (strong collision resistant)	It is computationally infeasible to find any pair $(x, y)$ with $x \neq y$ , such that $H(x) = H(y)$ .
Pseudorandomness	Output of $H$ meets standard tests for pseudorandomness.



# Secure Hash Algorithm (SHA)

- SHA was developed by NIST - 1993

Algorithm	Message Size	Block Size	Word Size	Message Digest Size
SHA-1	$< 2^{64}$	512	32	160
SHA-224	$< 2^{64}$	512	32	224
SHA-256	$< 2^{64}$	512	32	256
SHA-384	$< 2^{128}$	1024	64	384
SHA-512	$< 2^{128}$	1024	64	512
SHA-512/224	$< 2^{128}$	1024	64	224
SHA-512/256	$< 2^{128}$	1024	64	256

*Note:* All sizes are measured in bits.

# Applications Of Hash Functions

- Message authentication is achieved using a message authentication code (MAC).
- File Integrity check.

# File integrity check

## Only the first few images are available! Where are the rest?

We don't store/serve the full set of ISO images for all architectures, to reduce the amount of space taken u

## Non-free Firmware

This Debian image build only includes Free Software where possible. However, many systems include ha those firmware files for those cases. See the Debian Wiki [non-free firmware](#) page for more information.

## Other questions?

See the Debian CD [FAQ](#) for lots more information about Debian CDs and installation.

The images here were put together by the [Debian CD team](#) , using debian-cd and other software.

<a href="#">Name</a>	
<a href="#">Parent Directory</a>	
	<a href="#">SHA256SUMS</a>
	<a href="#">SHA256SUMS.sign</a>
	<a href="#">SHA512SUMS</a>
	<a href="#">SHA512SUMS.sign</a>
	<a href="#">debian-12.9.0-amd64-netinst.iso</a>
	<a href="#">debian-edu-12.9.0-amd64-netinst.iso</a>
	<a href="#">debian-mac-12.9.0-amd64-netinst.iso</a>

Apache/2.4.58 (Unix) Server at cdimage.debian.org Port 443

# Message authentication code (MAC)

- A function of the message and a secret key that produces a fixed-length value that serves as the authenticator.

$$\text{MAC} = C(K, M)$$

where

$M$  = input message

$C$  = MAC function

$K$  = shared secret key

MAC = message authentication code

# Message Authentication Code - MAC

- The message plus MAC are transmitted to the intended recipient.
- The recipient performs the same calculation on the received message, using the same secret key, to generate a new MAC.
- The received MAC is compared to the calculated MAC.

# MACs BASED ON HASH FUNCTIONS: HMAC

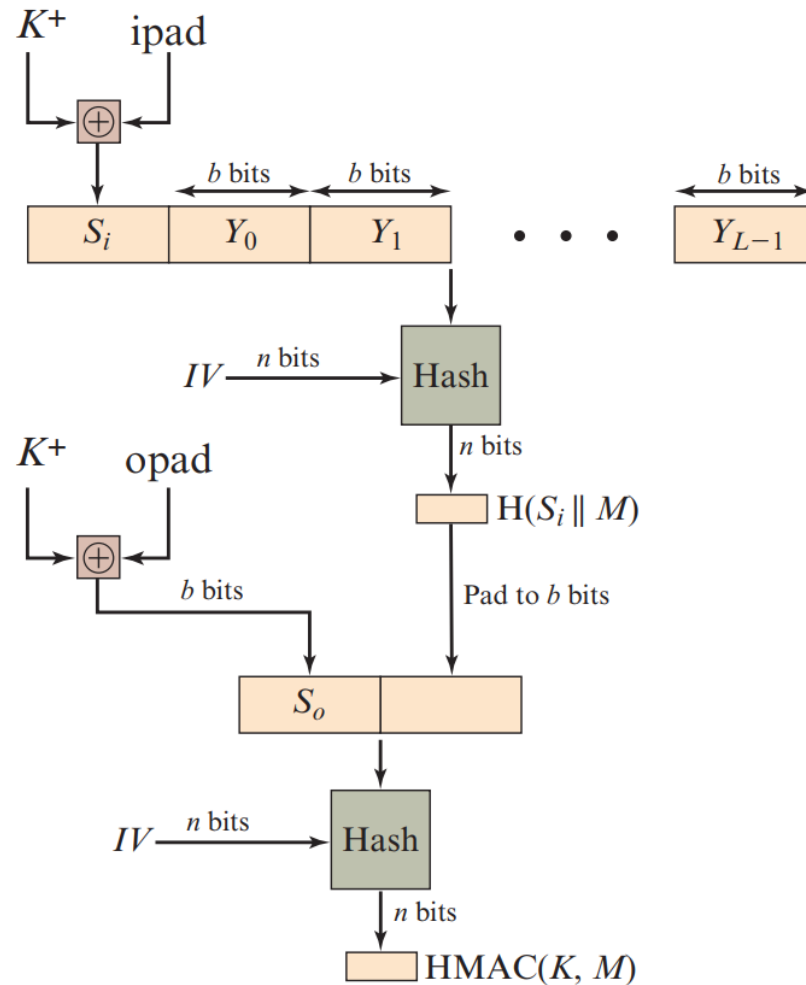


Figure 12.5 HMAC Structure